

1 Let  $(D, \sqsubseteq)$  be a poset. Prove that least upper bound is unique; i.e., prove the following theorem.

**Theorem.**  $\forall M \subseteq D, \forall x, y \in D, x = \sqcup M \wedge y = \sqcup M \implies x = y.$

In problems 2–3, let  $\mathbf{N}^+ = \{1, 2, 3, \dots\}$ , and  $\sqsubseteq$  be the binary relation on  $\mathbf{N}^+$  defined by  $x \sqsubseteq y$  if  $x$  divides evenly into  $y$  (with no remainder). For example,  $3 \sqsubseteq 12$  and  $3 \not\sqsubseteq 10$ . Prove the following theorems.

2 **Theorem.**  $(\mathbf{N}^+, \sqsubseteq)$  is a CPO.

3 Let  $f: \mathbf{N}^+ \rightarrow \mathbf{N}^+$  be defined by

$$f(x) = \begin{cases} x, & \text{if } x \leq 100 \\ 2x, & \text{if } x > 100. \end{cases}$$

**Theorem.**  $f$  is monotone  $\wedge f$  is not continuous.

4 Consider the following context-free grammar:

- Syntax

$$\begin{array}{lll} G & \longrightarrow & L \mid L\Delta \quad (\text{Grade}) \\ L & \longrightarrow & A \mid B \mid C \mid D \mid F \quad (\text{Letter}) \\ \Delta & \longrightarrow & + \mid - \quad (\text{Plus-Minus}) \end{array}$$

Design a denotational semantics specification in which  $\mathbf{R}_\perp$  is the flat CPO of real numbers, and  $\mathcal{G}: \text{Grade} \rightarrow \mathbf{R}_\perp$  is defined such that

$$\mathcal{G}[G] = \begin{cases} 4.0, & \text{if } G \text{ is } A \\ 3.7, & \text{if } G \text{ is } A- \\ 3.3, & \text{if } G \text{ is } B+ \\ 3.0, & \text{if } G \text{ is } B \\ 2.7, & \text{if } G \text{ is } B- \\ 2.3, & \text{if } G \text{ is } C+ \\ 2.0, & \text{if } G \text{ is } C \\ 1.7, & \text{if } G \text{ is } C- \\ 1.3, & \text{if } G \text{ is } D+ \\ 1.0, & \text{if } G \text{ is } D \\ 0.7, & \text{if } G \text{ is } D- \\ 0.0, & \text{if } G \text{ is } F \\ \perp, & \text{otherwise.} \end{cases}$$

5  $(D, \sqsubseteq)$  is a *lattice* if

- 1)  $(D, \sqsubseteq)$  is a partially-ordered set
- 2)  $\forall M \subseteq D, \exists x \in D, x = \sqcup M$
- 3)  $\exists$  least element  $\perp_D \in D$ .

Prove the following theorems:

- a) **Theorem.**  $\forall (D, \sqsubseteq), (D, \sqsubseteq)$  is a lattice  $\implies (D, \sqsubseteq)$  is a CPO.
- b) **Theorem.**  $\exists (D, \sqsubseteq), (D, \sqsubseteq)$  is a CPO  $\wedge (D, \sqsubseteq)$  is not a lattice.

6 Consider the following denotational semantics specification:

- Syntax

$$\begin{array}{ll}
 R & \longrightarrow L & \text{(Roman-Numeral)} \\
 L & \longrightarrow \varepsilon \mid LT & \text{(List)} \\
 T & \longrightarrow D \mid D_1 D_2 & \text{(Term)} \\
 D & \longrightarrow \text{I} \mid \text{V} \mid \text{X} \mid \text{C} \mid \text{M} & \text{(Digit)}
 \end{array}$$

- Semantic Functions

$$\begin{array}{l}
 \mathcal{R}: \text{Roman-Numeral} \rightarrow \mathbf{N}_\perp \\
 \mathcal{L}: \text{List} \rightarrow \mathbf{N}_\perp \times \mathbf{N}_\perp \\
 \mathcal{T}: \text{Term} \rightarrow \mathbf{N}_\perp \\
 \mathcal{D}: \text{Digit} \rightarrow \mathbf{N}_\perp
 \end{array}$$

are defined by

$$\begin{array}{l}
 \mathcal{R}[L] = \text{on}_1(\mathcal{L}[L]) \\
 \mathcal{L}[\varepsilon] = (0, 1000) \\
 \mathcal{L}[LT] = (\text{if } <(\text{smash}(\text{on}_2(\mathcal{L}[L]), \mathcal{T}[T])) \text{ then } \perp \text{ else } \text{sum}(\text{on}_1(\mathcal{L}[L]), \mathcal{T}[T]), \mathcal{T}[T]) \\
 \mathcal{T}[D] = \mathcal{D}[D] \\
 \mathcal{T}[D_1 D_2] = \text{if } <(\text{smash}(\mathcal{D}[D_1], \mathcal{D}[D_2])) \text{ then } \text{diff}(\text{smash}(\mathcal{D}[D_2], \mathcal{D}[D_1])) \text{ else } \perp \\
 \mathcal{D}[\text{I}] = 1 \\
 \mathcal{D}[\text{V}] = 5 \\
 \mathcal{D}[\text{X}] = 10 \\
 \mathcal{D}[\text{C}] = 100 \\
 \mathcal{D}[\text{M}] = 1000
 \end{array}$$

- a) Show that the context-free grammar is ambiguous, by giving two distinct derivation trees for the string of terminals **MCM**.
- b) Compute  $\mathcal{R}[\text{MCM}]$  separately for each of the derivation trees in a).