

Homework 3

CS 531
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For each of the following decision problems,

- Give a reasonable length function,
- Give a pseudo-code algorithm which correctly accepts instances for which the answer to the question is yes, and
- Give an upper bound on the worst-case time or space complexity of your algorithm—as a function of the length of an instance (defined by the length function)—to justify your result.

1 Show that **TSP** \in NP.

Traveling Salesman (TSP)

Instance: A finite set $C = \{c_1, c_2, \dots, c_m\}$ of “cities”,
a “distance” $d(c_i, c_j) \in \mathbb{Z}^+$ for each pair of cities $c_i, c_j \in C$, and
a bound $B \in \mathbb{Z}^+$ (where \mathbb{Z}^+ denotes the positive integers).

Question: Is there a “tour” of all the cities in C having total length no more than B , i.e.,
an ordering $\langle c_{\pi(1)}, c_{\pi(2)}, \dots, c_{\pi(m)} \rangle$ of C such that

$$\left(\sum_{i=1}^{m-1} d(c_{\pi(i)}, c_{\pi(i+1)}) \right) + d(c_{\pi(m)}, c_{\pi(1)}) \leq B?$$

2 Show that **SAT** \in NP.

Satisfiability (SAT)

Instance: A set U of variables and a collection C of clauses over U .

Question: Is there a satisfying truth assignment for C ?

3 Show that **QBF** \in PSPACE.

Quantified Boolean Formulas (QBF)

Instance: A well-formed quantified Boolean formula $F = (Q_1x_1)(Q_2x_2) \cdots (Q_nx_n)E$,
where E is a Boolean expression involving the variables x_1, x_2, \dots, x_n and
each Q_i is either “ \exists ” or “ \forall ”.

Question: Is F true?

4 Show that **GAP** \in NL.

Graph Accessibility Problem (GAP)

Instance: An undirected graph $G = (V, E)$, and vertices $a, b \in V$.

Question: Is there a path in G from a to b ?